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ON COMBINING PSEUDORANDOM NUMBER GENERATORS

by

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Introduction

Many methods have been proposed, tested and employed for generating pseudorandom numbers ([2], [3], [4], [5], [7], [9], [10], [12], [13]). The goal is to produce strings of numbers which behave like independent uniform $[0,1]$ random variables. The generators yield integers in the set $\{0, \dots, m-1\}$ which are then transformed to $[0,1]$ by division by m .

Suppose that X_1, X_2, \dots and Y_1, Y_2, \dots are strings of numbers in $\{0, \dots, m-1\}$ generated by two separate generators. Assume that the two strings are independent. Define a new string of numbers Z_1, Z_2, \dots by $Z_i = X_i + Y_i \pmod{m}$. For any k and corresponding i_1, \dots, i_k define $\underline{X} = (X_{i_1}, \dots, X_{i_k})$, $\underline{Y} = (Y_{i_1}, \dots, Y_{i_k})$, $\underline{Z} = (Z_{i_1}, \dots, Z_{i_k})$. Let r be the distribution of k independent random variables uniformly distributed on $\{0, \dots, m-1\}$. We consider several natural measures of distance between multivariate distribution with components in $\{0, \dots, m-1\}$ and show under these distances that Z_{i_1}, \dots, Z_{i_k} has a distribution closer to r than either X_{i_1}, \dots, X_{i_k} or Y_{i_1}, \dots, Y_{i_k} for all k, i_1, \dots, i_k .

In applying our results to pseudorandom number generation, two points need careful scrutiny. First, we assume that the two strings, X and Y , are independent. Secondly, X, Y and Z are deterministic, even though they are being constructed to look random. This creates a problem in the interpretation of Lemma 1 and in the interpretation of independence. We certainly do not claim that our results prove that addition mod (m) of separately generated sequences improves pseudorandom number generation. We only assert that this conclusion is suggested and warrants further study.

The technique of combining strings by addition mod m is also mentioned in Knuth ([5], p. 30). An exercise at the end of the section (p. 33) shows that if the periods of \underline{X} and \underline{Y} are λ_1, λ_2 with λ_1 and λ_2 relatively prime, then the period of \underline{Z} is $\lambda_1 \lambda_2$. This implies that we should choose the periods of the separate generators to be relatively prime.

Random number generators are generally studied by statistical tests on the output, and by mathematical analysis of the period. We hope that the methods employed here will provide another approach to analyzing pseudorandom number generators.

Results

Suppose that $\underline{X} = (X_1, \dots, X_k)$ and $\underline{Y} = (Y_1, \dots, Y_k)$ are independent random vectors with $\Pr(X_1=j_1, X_2=j_2, \dots, X_k=j_k) = p(j_1, \dots, j_k)$, $\Pr(Y_1=l_1, Y_2=l_2, \dots, Y_k=l_k) = q(l_1, \dots, l_k)$; each component assumes values in $\{0, \dots, m-1\}$. Define $\underline{Z} = (Z_1, \dots, Z_k)$ with $Z_i = X_i + Y_i \pmod{m}$, and $\Pr(Z_1=m_1, \dots, Z_k=m_k) = s(m_1, \dots, m_k)$. As measures of departure of a distribution b on $\prod_{i=1}^k \{0, \dots, m-1\}_i$, from r the distribution of k independent uniforms on $\{0, \dots, m-1\}$, we use:

$$(i) \quad \|b - r\|_\alpha = \sum_{(j_1, \dots, j_k)} \left| b(j_1, \dots, j_k) - \frac{1}{m^k} \right|^\alpha, \quad 1 \leq \alpha < \infty,$$

$$(ii) \quad \|b - r\|_\infty = \max_{(j_1, \dots, j_k)} \left| b(j_1, \dots, j_k) - \frac{1}{m^k} \right|,$$

$$(iii) \quad \Pi(b, r) = \sum_{(j_1, \dots, j_k)} b(j_1, \dots, j_k) \log(m^k b(j_1, \dots, j_k)).$$

Quantity (iii) is the mean information in favor of b against the distribution r (see Kullback [6], p. 5); $\Pi(b, r)$ achieves a minimum of 0 at $b = r$ and is otherwise positive.

Lemma 1: For $1 \leq \alpha \leq \infty$, $\|s - r\|_\alpha \leq \min(\|p - r\|_\alpha, \|q - r\|_\alpha)$, and $\Pi(s, r) \leq \min(\Pi(p, r), \Pi(q, r))$.

Proof of Lemma 1: We rely heavily on the technique of majorization ([1], [8], [11]). Firstly, adding \underline{Y} to an independent random variable \underline{X} is equivalent to making a transition in the Markov chain with transition matrix $P_{(j_1, \dots, j_k), (l_1, \dots, l_k)} = r(l_1 - j_1 \pmod{m}, l_2 - j_2 \pmod{m}, \dots, l_k - j_k \pmod{m})$. Now $\sum_{(j_1, \dots, j_k)} P_{(j_1, \dots, j_k), (l_1, \dots, l_k)} = \sum_{(m_1, \dots, m_k)} r(m_1, \dots, m_k) = 1$ ((m_1, \dots, m_k) ranges through the m^k sample points achieving each exactly once). Therefore P is doubly stochastic. Next, $s = pP$ with P doubly stochastic and it thus follows from a theorem of Karamata ([1], pg. 31) that s is majorized by p . By definition this means that if s and p are rearranged so that $s_{(1)} \geq s_{(2)} \geq \dots \geq s_{(m^k)}$ and

$$p_{(1)} \geq p_{(2)} \geq \dots \geq p_{(m^k)}, \text{ then } \sum_{i=1}^j s_{(i)} \leq \sum_{i=1}^j p_{(i)} \text{ for}$$

$$j = 1, 2, \dots, m^k - 1, \text{ and } \sum_{i=1}^{m^k} s_{(i)} = \sum_{i=1}^{m^k} p_{(i)}.$$

It follows from the definition of majorization that if s is majorized by p then $s - \frac{1}{m^k}$ is majorized by $p - \frac{1}{m^k}$. For $1 \leq \alpha \leq \infty$, $|x|^\alpha$ is continuous and convex; it then follows from [1], p. 30, that

$$\|s - r\|_\alpha \leq \|p - r\|_\alpha.$$

From the definition of majorization, we know that $s_{(1)} \leq p_{(1)}$ while $s_{(m^k)} \geq p_{(m^k)}$. Therefore

$$\|s - r\|_{\infty} = \max(s_{(1)} - \frac{1}{m^k}, \frac{1}{m^k} - s_{(m^k)}) \leq \max(p_{(1)} - \frac{1}{m^k}, \frac{1}{m^k} - p_{(m^k)}) = \|p - r\|_{\infty}.$$

Finally, the function $F(x_1, \dots, x_{m^k}) = \sum_{i=1}^{m^k} x_i \log(m^k x_i)$ satisfies $(x_i - x_j) \left(\frac{\partial F}{\partial x_i} - \frac{\partial F}{\partial x_j} \right) = (x_i - x_j) \log \frac{x_i}{x_j} \geq 0$ for $x_i \geq x_j$. It then follows from a theorem of Ostrowski ([1], p. 32) and the majorization of s by p that $\Pi(s, r) \leq \Pi(p, r)$.

By reversing the roles of X and Y it follows that $\|s - r\|_{\alpha} \leq \|q - r\|_{\alpha}$, $1 \leq \alpha \leq \infty$, and $\Pi(s, r) \leq \Pi(q, r)$. This concludes the proof.

Undoubtedly our result will hold for many other metrics.

If we take an independent sequence of random vectors X_1, \dots, X_n, \dots and form partial sums (mod m), $Z_n = \sum_{i=1}^n X_i \pmod{m}$, $n = 1, 2, \dots$, we will, under weak conditions, converge at a geometric rate to r . This follows from standard Markov chain analysis.

More specifically, assume that

$$\min_{(j_1, \dots, j_k)} \Pr(X_{(i),1} = j_1, \dots, X_{(i),k} = j_k) = \Delta_i \geq \Delta > 0$$

for all i . Then, letting s_n denote the distribution of Z_n , it follows that

$$\begin{aligned} \max_{(j_1, \dots, j_k)} s_n(j_1, \dots, j_k) - \min_{(j_1, \dots, j_k)} s_n(j_1, \dots, j_k) &\leq \sum_{i=1}^n (1 - m^k \Delta_i) \\ &\leq (1 - m^k \Delta)^n \rightarrow 0 \end{aligned}$$

as $n \rightarrow \infty$. For interpreting this result, note that

$$\begin{aligned} & \max_{(j_1, \dots, j_k)} s_n(j_1, \dots, j_k) - \min_{(j_1, \dots, j_k)} s_n(j_1, \dots, j_k) \\ &= \max_{(j_1, \dots, j_k), (\ell_1, \dots, \ell_k)} \left| s_n(j_1, \dots, j_k) - s_n(\ell_1, \dots, \ell_k) \right| \\ &\geq \max_{(j_1, \dots, j_k)} \left| s_n(j_1, \dots, j_k) - \frac{1}{m^k} \right|. \end{aligned}$$

The proof is simple. Let $M_n = \max_{(j_1, \dots, j_k)} s_n(j_1, \dots, j_k)$ and $m_n = \min_{(j_1, \dots, j_k)} s_n(j_1, \dots, j_k)$. Then

$$M_n \leq M_{n-1}(1 - (m^k - 1)\Delta_n) + \Delta_n(1 - M_{n-1}) = M_{n-1}(1 - m^k \Delta_n) + \Delta_n,$$

while

$$m_n \geq m_{n-1}(1 - (m^k - 1)\Delta_n) + \Delta_n(1 - m_{n-1}) = m_{n-1}(1 - m^k \Delta_n) + \Delta_n.$$

Thus $M_n - m_n \leq (M_{n-1} - m_{n-1})(1 - m^k \Delta_n)$. Repeated use of this argument gives

$$(M_n - m_n) \leq (M_1 - m_1) \prod_{i=2}^n (1 - m^k \Delta_i) \leq \sum_{i=1}^n (1 - m^k \Delta_i) \leq (1 - m^k \Delta)^n.$$

Under the weaker condition $\sum_{i=1}^{\infty} \Delta_i = \infty$, we still get

$\lim_{n \rightarrow \infty} (M_n - m_n) = 0$, although not necessarily convergence at a geometric rate. The condition $\sum_{i=1}^{\infty} \Delta_i = \infty$ is sufficient but not necessary for

convergence of $M_n - m_n$ to 0 . For example, if any one $\Delta_i = \frac{1}{m^k}$
(equivalently if any one $X_{(i)}$ has distribution r) then the distribu-
tion of Z_n is r for all $n \geq i$.

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20. Let $\underline{X} = (X_1, \dots, X_n)$ and $\underline{Y} = (Y_1, \dots, Y_n)$ be independent random vectors whose components take values in $\{0, 1, \dots, m-1\}$. Let r be the joint distribution of n independent random variables uniformly distributed on $\{0, 1, \dots, m-1\}$. We show that the distribution of $\underline{Z} = \underline{X} + \underline{Y} \pmod{m}$ is closer to r , in several metrics, than is either the distribution of \underline{X} or of \underline{Y} . The principle suggested by this result is that combining strings of pseudorandom numbers, generated by different generators, by addition mod m , will result in a string more random than any of the separate strings.

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